### **Fermat's Last Theorem**

**Diophantine equation** 

$$x^n + y^n = z^n$$

does not have a solution in positive integers for integer n > 2.

- Conjectured by Pierre de Fermat in 1637 in the margin of a copy of Arithmetica where he claimed he had a proof that was too large to fit in the margin.
- $\bigcirc$  *n* = 4 Fermat + several equivalent proofs
- n = 3 Leonhard Euler (1770)
- n = 5 Legendre / Dirichlet (1825)
- **o** *n* = 7 Lamé (1839)

● *n* = 6, 10, 14

- ł
- any *n* Andrew Wiles (1994)
- Elliptic curves  $y^2 = x^3 + ax + b$
- Modular forms  $\mathbb{C}^2 \to \mathbb{C}^2$  with "much symmetry"

n = 2 (Euclid's formula):  $x = m^2 - n^2$  y = 2mn $z = m^2 + n^2$ 

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#### **Fermat's Little Theorem**

For p = prime:

gcd = greatest common divisor

 $a^p \equiv a \pmod{p}$   $a^{p-1} \equiv 1 \pmod{p}, a \text{ not multiple of } p$ 

**Proof:** Consider *p*-tuples of *a* objects; there are  $a^p$  of them. We remove 111..1, 222..2,...; there are  $a^p - a$  left. These can be grouped to *p*-cyclically shifted groups; e.g., 21111, 12111, 11211, 11121, 11121, 11112.

Also:  $\{a, a^2, a^3, \dots, a^{p-1}\} = \{1, 2, 3, \dots, p-1\} \pmod{p}$  (i.e., except order)

**Extension:** for *a*, *n* co-primes

$$a^{\phi(n)} \equiv 1 \pmod{n}$$

numbers a, b so that gcd(a, b) = 1are called co-primes

where  $\phi(n) = \text{Euler's totient function} = \text{number of co-primes to } n \text{ in interval } [1, n-1].$ NB:  $\phi(p) = p - 1.$ 

**Example:** calculate 3<sup>7</sup> (mod 7) by the square-and-multiply algorithm

 $(\text{mod } 7): 3^2 \equiv 2, 3^4 \equiv 4, 3^6 \equiv 1, 3^7 \equiv 3$ 

**Inversion:** If  $a^{n-1} \not\equiv 1$  for a co-prime a, then n is composite

**Probabilistic test:** If  $a^{n-1} \equiv 1$  for several co-primes a, then n is a prime with a high probability

### **Modular inversion (division in general)**

Let a, b are co-primes and a < b. We want to solve

 $ax \equiv 1 \pmod{b}$  or ax + by = 1

#### **Extended Euclidean algorithm:**

$$a \qquad b$$

$$r_{0} := b \qquad s_{0} = 0 \qquad t_{0} = 1$$

$$r_{1} := a \qquad s_{1} = 1 \qquad t_{1} = 0$$

$$r_{2} := r_{0} - q_{1}r_{1} \qquad s_{2} := s_{0} - q_{1}s_{1} \qquad t_{2} := t_{0} - q_{1}t_{1}$$

$$r_{3} := r_{1} - q_{2}r_{2} \qquad s_{3} := s_{1} - q_{2}s_{2} \qquad t_{3} := t_{1} - q_{2}t_{2}$$

$$\vdots$$

$$1 \qquad x \qquad y$$

where  $q_i$  = reminder after  $r_{i-1}$  :  $r_i$  (: = integer division)

**Proof:** based on  $r_i = as_i + bt_i$  for every line, proof by induction.

**Example:** solve  $6x \equiv 1 \pmod{17}$ 

## **RSA cryptosystem**

Rivest–Shamir–Adleman (1978)

- Choose 2 distinct primes *p*, *q* (not too close)
- Calculate n = pq (the modulo, part of the public key)
- Calculate  $\lambda = (p-1)(q-1)$  (better: lcm(p-1, q-1))
- Public key: e,  $1 < e < \lambda$ , co-prime to  $\lambda$  (often e = 65537)
- Private key: d so that  $de \equiv 1 \pmod{\lambda}$  (easy if you know pq)

**Encrypt**  $m: c \equiv m^e \pmod{n}$ 

**Decrypt**  $c: c^d \equiv m \pmod{n}$ 

**Proof.** Since de - 1 is a multiple of  $\lambda$ ,  $\exists g, h, k$  so that

$$ed-1=g\lambda=h(p-1)=k(q-1)$$

Using Fermat's little theorem (except  $m \equiv 0 \pmod{p}$ , which is trivial)

$$m^{ed} = m^{ed-1}m = (m^{p-1})^h m \equiv 1^h m \equiv m \pmod{p}$$

And similarly for  $q: m^{ed} \equiv m \pmod{q}$ . Since p, q are co-primes,

 $(m^e)^d \equiv m \pmod{pq}$  q.e.d.

#### Icm = least common multiple

? Can integer factorization be solved in polynomial time on a classical computer?

### How it works

#### Message sent via insecure channel (https, ssh)

- Alice calculates n, e and sends it openly to Bob.
- Bob encrypts a message using *n*, *e* and sends it to Alice.
- Alice decrypts the mesage using her private n, d.

#### **Digital signature**

- Alice publishes n, e.
- Alice encrypts a file (better: a hash) using n, d.
- Bob can verify the encrypted hash using n, e.

#### **SSH login without password**

- Generate a private/public key pair on your HOME computer: ssh-keygen -t rsa your PRIVATE key is .ssh/id\_rsa your PUBLIC key is .ssh/id\_rsa.pub
- copy your PUBLIC key to .ssh/authorized\_keys on the REMOTE machine



start -/n pic/theories.jkv

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### **The Enormous Theorem**

Every finite simple group is isomorphic to one of the following groups:

- 1. A cyclic group with prime order;
- 2. An alternating group (group of even permutations) of degree at least 5;
- 3. A simple group of Lie type (over a finite field) (quite rich...)<sup> $\dagger$ </sup>
- 4. The 26 sporadic simple groups.

The biggest sporadic group = "Monster", number of elements

 $= 2^{46} \cdot 3^{20} \cdot 5^9 \cdot 7^6 \cdot 11^2 \cdot 13^3 \cdot 17 \cdot 19 \cdot 23 \cdot 29 \cdot 31 \cdot 41 \cdot 47 \cdot 59 \cdot 71$ 

= 80801742479451287588645990496171075700575436800000000

Proof finished 2004 – hundreds of papers...

Group is a set G with "multiplication" and "division":

- $\bigcirc \forall a, b \in G: ab \in G$  (closure)
- $\bigcirc$   $\forall a, b, c \in G: (ab)c = a(bc)$  (associativity)
- $\bigcirc$   $\exists e \in G : \forall a \in G \text{ it holds } ea = ae = a \text{ (identity element)}$
- $\bigcirc \forall a \in G \exists a^{-1} : aa^{-1} = a^{-1}a = e$  (inverse element)

<sup>†</sup> Incl. both the classical Lie groups, namely the simple groups related to the projective special linear, unitary, symplectic, or orthogonal transformations over a finite field; the exceptional and twisted groups of Lie type (incl. the Tits group, sometimes classified as sporadic).

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#### **Four color theorem**

Every map (on sphere or plane) can be colored by 4 colors



credit: http://www.artsrn.ualberta.ca/mengel/2015huco617/files/2015/04/Map.jpg

Computer-assisted proof in 1976 by Kenneth Appel and Wolfgang Haken, based on 1,936 submaps.

Easier for torus etc.

#### **Deterministic chaos**

Weather, oil on pan ... Lorentz attractor:

$$\dot{x} = \sigma y - \sigma x,$$
  

$$\dot{y} = \rho x - xz - y,$$
  

$$\dot{z} = xy - \beta z$$



cd ../maple; xmaple mmpc7.mw

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Simpler model:  $x := a - x^2$  (see mmpc7.mw)

universal properties; Feigenbaum: 4.669201609102990671853203821578... 2.502907875095892822283902873218...

self-similarity (fractal)

**Bifurcation map** x := rx(1-x)



#### **Fractals**

Problem: what is the length of the borderline? Answer: it depends on the meter *m*:

 $l = \operatorname{const} m^{1-D}$ 

D = 1.02 South Africa D = 1.25 west GB

**Fractal:** geometric set, which resembles a part of itself (after a continuous transformation, usually shrinking)

Random fractal: self-similarity in a statistical-sense

(Almost) definition of the **fractal dimension**:

$$D = \lim_{m \to 0} \frac{\log N_m}{\log(1/m)}$$

where  $N_m = \#$  of line segments/squares/cubes ... of length/edge... *m* needed to cover the set (1/m = # of line segments of length *m* to cover a unit line segment, D = 1)



show/fraktaly.sh

10/24

### **Fractal dimension**

**Example.** Calculate the fractal dimension of a line segment. Answer:  $N_m = l/m$ ,  $D = \lim \log(l/m) / \log(1/m) = 1$ 



**Example.** Calculate the fractal dimension of the Koch curve

10 4/ |u 3 ÷ 1,26

cd show;mz Kochsim.gif 11/24

mmpc7

**Example.** Calculate the fractal dimension of the trajectory of the Brownian motion (polymer in a  $\theta$ -solvent)

1 step of random walk by  $1 \begin{pmatrix} 1/2 \\ \leftarrow \end{pmatrix}$ ,  $\stackrel{1/2}{\rightarrow}$ ):  $\langle R^2 \rangle = 1$ , m = 1, l = 1,  $N_m = 1$ 

2 steps of random walk:  $\langle R^2 \rangle = 1$ ,  $m = 1/\sqrt{2}$ ,  $l = \sqrt{2}$ ,  $N_m = l/m = 2$ 

(noisnamib as rhe space dimension)  $\Delta = \Delta$ 

# **Poincaré hypothesis**

- Every simply-connected, closed 3-manifold is homeomorphic to the 3-sphere.
- 3-sphere =  $\{\vec{r}, |\vec{r}| = 1\}$  in  $\mathbb{R}^4$
- simply-connected = path-connected + any circle can be be contracted to a point
- $\bigcirc$  path-connected =  $\exists$  a continuouos path between points
- closed = compact + without boundary
- compact = any open cover has a finite subcover; any infinite sequence has a converging subsequence
- 3-manifold = locally as 3D Euclidean
- homeomorphism = continuous function between topological spaces that has a continuous inverse function
- Proven by Grigori Perelman 2002, 2003 He rejected the Millennium Prize (\$1M) and the Fields medal
- Cf. Poincaré homology sphere (glued dodekahedron, binary icosahedral group, n = 120)

# **The Rectangular Peg Problem**

aka Inscribed Square Problem, Toeplitz' conjecture

Does any (smooth) Jordan curve admit an inscribed square?

A Jordan curve is a plane curve which is topologically equivalent to (a homeomorphic image of) the unit circle, i.e., it is simple and closed.

Solved by Joshua Evan Greene and Andrew Lobb (2020) for smooth curves (even for every rectangle  $\exists$  a similar rectangle on the the curve)



credits: Wikipedia, https://www.quantamagazine.org/new-geometric-perspective-cracks-old-problem-about-rectangles-20200625/

mz movies/RectangularPegProblem.gif 13/24 mmpc7

#### Gödel

Liar paradox: This sentence is false.

**Paradox of 1000 words:** Let us consider sentences in English, based on a well-defined syntax and vocabulary, that define a number. There are a finite number of words in English; therefore, there are a finite number of sentences shorter than 1000 words that define a number. Hence, there exists a highest such number. Now, consider the sentence: "The number equals one plus the highest number that can be defined by a sentence in English composed of less than 1000 words."

Language vs. metalanguage

**First Incompleteness Theorem:** Any consistent formal system *F* within which a certain amount of elementary arithmetic can be carried out is incomplete; i.e., there are statements of the language of *F* which can neither be proved nor disproved in *F*.

**First Incompleteness Theorem:** For each formal system *F* containing basic arithmetic, it is possible to canonically define a formula Cons(F) expressing the consistency of *F*. This formula expresses the property that "there does not exist a natural number coding a formal derivation within the system *F* whose conclusion is a syntactic contradiction;" loosely: "there is no natural number that codes a derivation of 0=1 from the axioms of *F*."

Gödel's second incompleteness theorem shows that, under general assumptions, this canonical consistency statement Cons(F) will not be provable in F.

### **TO ADD: Axiom of choice**

### **Curves/shapes of constant width**

A **curve of constant width** is a simple closed curve in the plane whose width (the distance between parallel supporting lines) is the same in all directions. Such curve is convex.





cd show; firefox ReuleauxTriangleAnimation.gif 16/24

cd ../simul/kissing; kiss.sh 17/24

### **? Kissing numbers**

The greatest number of non-overlapping unit spheres (in Ddimensional space) touching a central unit sphere.



D	Lower bound	Upper bound	proven
1	2		easy
2	6		easy
3	12		1953
4	24		2003
5	40	44	
6	72	78	
7	126	134	
8	$240 = 2^4 \cdot 3 \cdot 5$		1979
9	306	364	
10	500	554	
11	582	870	
12	840	1,357	
13	1,154	2,069	
14	1,606	3,183	
15	2,564	4,866	
16	4,320	7,355	
17	5,346	11,072	
18	7,398	16,572	
19	10,668	24,812	
20	17,400	36,764	
21	27,720	54,584	
22	49,896	82,340	
23	93,150	124,416	
24	$196,560 = 2^4 \cdot 3^3 \cdot 5 \cdot 7 \cdot 13$		1979

# **? Twin primes**

are prime numbers  $p_1, p_2$  so that  $p_2 - p_1 = 2$ .

Are there infinitely many twin primes?

Probably yes, but not proven...

Brun's theorem: the sum of reciprocal twin primes converges:

$$\sum_{p,p+2 \text{ are primes}} \left(\frac{1}{p} + \frac{1}{p+2}\right) = \left(\frac{1}{3} + \frac{1}{5}\right) + \left(\frac{1}{5} + \frac{1}{7}\right) + \left(\frac{1}{11} + \frac{1}{13}\right) + \dots \approx 1.902160583104$$

Euler 1737: the sum of reciprocal primes diverges

Yitang Zhang 2015:  $\liminf_{n\to\infty} (p_{n-1}-p_n) < 246$ 

### **? Goldbach's conjecture**

Every even integer greater than 2 can be expressed as a sum of two primes.

The conjecture has been shown to hold for all integers less than  $4 \times 10^{18}$ 

### **? Odd perfect number**

**Perfect number** = positive integer that is equal to the sum of its positive divisors, excluding the number itself.

Euclid proved that  $2^{p-1}(2^p-1)$  is an even perfect number whenever  $2^p-1$  is prime (Mersenne prime).

$$6 = 110_2$$
  $28 = 11100_2$   $496 = 111110000_2$ 

It is unknown whether there is any odd perfect number *N*. If yes:

 $\bigcirc N > 10^{1500}$ 

The largest prime factor of N is greater than  $10^8$ 

N has at least 101 prime factors and at least 10 distinct prime factors.

### **Catalan's conjecture** → **Mihăilescu's theorem**

The only solution in natural numbers of  $x^a - y^b = 1$  for a, b > 1, x, y > 0 is  $3^2 - 2^3 = 1$ .

Proven 2002 by Preda V. Mihăilescu

#### ? Zarankiewicz and Hill conjectures



To fully connect, algorithms exist for the number of crossing lines C:

$$C(K_{n,m}) = \lfloor \frac{n}{2} \rfloor \lfloor \frac{n-1}{2} \rfloor \lfloor \frac{m}{2} \rfloor \lfloor \frac{m-1}{2} \rfloor \quad C(K_n) = \frac{1}{4} \lfloor \frac{n}{2} \rfloor \lfloor \frac{n-1}{2} \rfloor \lfloor \frac{n-2}{2} \rfloor \lfloor \frac{n-3}{2} \rfloor$$

Are there better solutions?

[M. Balko, Vesmír **11**, 628 (2019)]

## **?** Collatz conjecture aka 3n + 1 problem

 $f(n) := \begin{cases} n/2 & \text{for even } n \\ 3n+1 & \text{for odd } n \end{cases}$ 

**Conjecture:** For any integer *n*, there exists a finite *M* 

so that 
$$f(f(\ldots(n))) = 1$$

• proven for  $n < 2^{68}$ 

- probabilistic arguments suggest it is true
- The cycle length (except  $1 \rightarrow 2 \rightarrow 1...$ ) of

$$f(n) := \begin{cases} n/2 & \text{for even } n \\ (3n+1)/2 & \text{for odd } n \end{cases}$$

is at least 17087915 (Eliahou 1993). Based on continuum fraction expansion for  $\ln 3/\ln 2 \rightarrow \rightarrow \rightarrow$ 

27 82 41 124 62 31 94 47 142 71 214 107 322 161 484 242 121 364 182 91 274 137 412 206 103 310 155 466 233 700 350 175 526 263 790 395 1186 593 1780 890 445 1336 668 334 167 502 251 754 377 1132 566 283 850 425 1276 638 319 958 479 1438 719 2158 1079 3238 1619 4858 2429 7288 3644 1822 911 2734 1367 4102 2051 6154 3077 9232 4616 2308 1154 577 1732 866 433 1300 650 325 976 488 244 122 61 184 92 46 23 70 35 106 53 160 80 40 20 10 5 16 8 4 2 1 4 2 1 4 2 1

1	1/1	-0.585
1	2/1	0.415
1	3/2	-0.08496
2	8/5	0.01504
2	19/12	-0.001629
3	65/41	0.0004034
1	84/53	-5.684e-05
5	485/306	4.82e-06
2	1054/665	-9.471e-08
23	24727/15601	1.683e-09
2	50508/31867	-3.289e-10
2	125743/79335	6.664e-11
1	176251/111202	-4.671e-11
1	301994/190537	4.883e-13
55	16785921/10590737	-7.243e-15
1	17087915/10781274	1.515e-15

# ? Riemann hypothesis

Riemann zeta function:

$$\zeta(s) = \sum_{n=1}^{\infty} \frac{1}{n^s} = \frac{1}{1^s} + \frac{1}{2^s} + \frac{1}{3^s} + \cdots$$

for  $s \in \mathbb{C}$  and  $\Re(s) > 1$ , and its analytic continuation

Euler:

$$\zeta(s) = \prod_{p \text{ is prime}} \frac{1}{1 - p^{-s}}$$

Single pole at s = 1 (res = 1)

Hypothesis (1859): roots = negative even integers (trivial) and complex numbers with real part 1/2.

Proven (2004) for the first  $10^{13}$  roots, but not in general

Consequences to the distribution of primes



credit: http://wismuth.com/complex/gallery.html

# **? Complexity theory:** *P* = *NP*

P = problem can be solved (on a computer) in a polynomial time (as a function of problem size)\*
 e.g.: sorting, square root
 Subset sum problem: Having data

 $NP^{\dagger}$  = a known solution can be verified in a polynomial time e.g., subset sum problem, sudoku

Subset sum problem: Having data set (numbers may repeat) of *N* integers and target sum *T*, find a data subset summing up to *T* 

**NP-complete** = problems to which any other NP-problem can be reduced in polynomial time, and whose solution may still be verified in polynomial time e.g., decide whether a solution of traveling salesman is indeed the shortest

NP-hard = at least as hard as the hardest NP H is NP-hard if every NP problem can be reduced in polynomial time to H e.g., traveling salesman, quantum theory, ...

● likely but not proven: P ≠ NP ⇒ NP-hard problems cannot be solved in polynomial time

\*for number problems: size = # of digits
<sup>†</sup>Nondeterministic Polynomial

#### **Connes embedding conjecture**

(about aproximating operators ≈
infinite matrices ≈ quantum world
by a hirearachy of finite matrices)
has been falsified

# **? Probably** $P \neq NP$

.. but not proven! And thus many problems are hard.



credit: "P np np-complete np-hard" by Behnam Esfahbod. Licensed under CC BY-SA 3.0 via Commons – https://commons.wikimedia.org/wiki/File:P\_np\_np-complete\_np-hard.svg#/media/File:P\_np\_np-cc